

AD-A198 366

TUTORIAL TRACK II ADVANCED ADA TOPICS(U) INFORMATION
SYSTEMS AND TECHNOLOGY CENTER W-P AFB OH ADA VALIDATION
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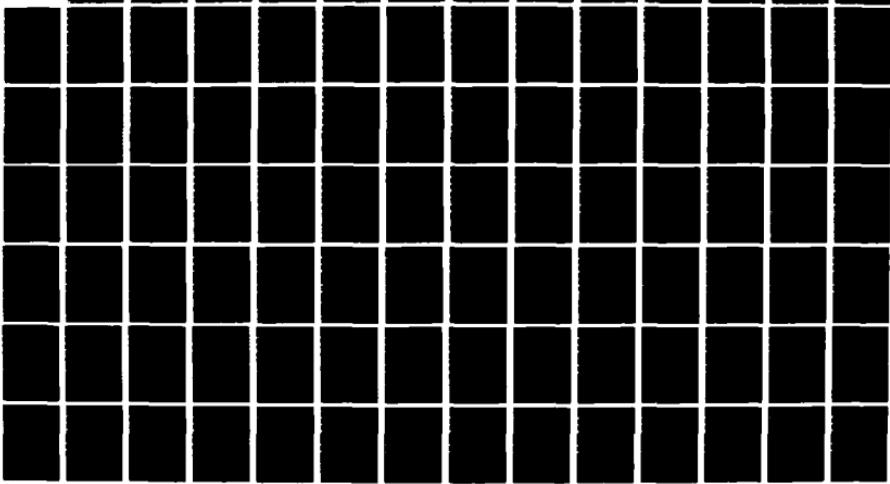
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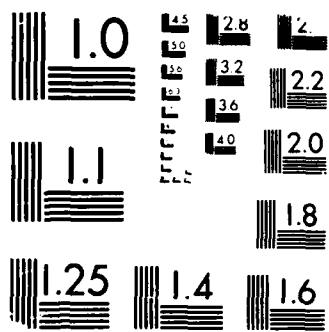
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AD-A190 366

TUTORIAL

TRACK II

ADVANCED ADA TOPICS

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Abstraction, Tasking, Strong Typing, and Exceptions

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Ada Tasking*

Abstraction of Process

by

Dean W. Gonzalez

David A. Cook

303-472-2136
AV 259-2136

U.S. Air Force Academy

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Ada Joint Program Office.

ADA TASKING

- OVERVIEW

DEFINE ADA TASKING

DEFINE SYNCHRONIZATION MECHANISM

EXAMPLES



- STARTED AFTER ELABORATION OF PARENT, AND BEFORE THE PARENT'S FIRST STATEMENT
- MAY ALSO BE A TYPE AND TREATED AS AN OBJECT

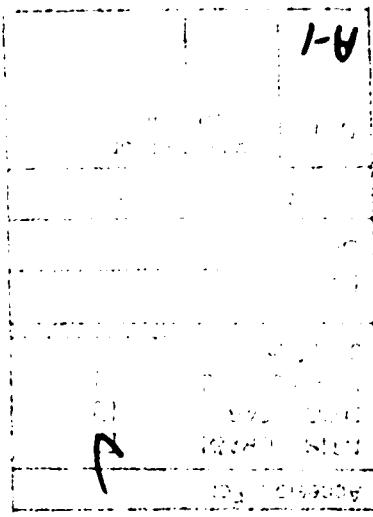
ADA TASKING

TASK DEFINITION

- A PROGRAM UNIT FOR CONCURRENT EXECUTION

- NEVER A LIBRARY UNIT

- MASTER IS A ...
 - LIBRARY PACKAGE
 - SUBPROGRAM
 - BLOCK STATEMENT
 - OTHER TASK



CALLEE PROVIDES SERVICE

1. IMMEDIATE RESPONSE
2. WAIT FOR A WHILE
3. WAIT FOREVER

SERVICE IS REQUESTED WITH AN ENTRY
CALL STATEMENT

SERVICE IS PROVIDED WITH AN ACCEPT
STATEMENT



ADA TASKING

SYNCHRONIZATION MECHANISMS

- GLOBAL VARIABLES

- RENDEZVOUS

- MAIN PROGRAM IN A TASK

- CALLER REQUESTS SERVICE

1. IMMEDIATE REQUEST
2. WAIT FOR A WHILE
3. WAIT FOREVER



ADA TASKING

SCENARIO I

"THE GOLDEN ARCHES"

MCD TASKS :

| | | |
|-------------------|---|------|
| SERVICE PROVIDED | : | Food |
| SERVICE REQUESTED | : | None |

GONZO TASKS :

| | | |
|-------------------|---|------|
| SERVICE PROVIDED | : | None |
| SERVICE REQUESTED | : | Food |



ADA TASKING

- SELECT STATEMENT PROVIDES ABILITY TO PROGRAM THE DIFFERENT 'REQUEST' AND 'PROVIDE' MODES
- GUARDS ARE "IF STATEMENTS" FOR THE PROVIDING SERVICE
- TERMINATION IS AN ALTERNATIVE IF A SERVICE IS NO LONGER NEEDED

Task McD is
entry SERVE(TRAY_OF : out FOOD_TYPE);
end McD;

Task GONZO;

Task Body McD is
NEW_TRAY : FOOD_TYPE;
function COOK return FOOD_TYPE is.....
begin
loop
accept SERVE (TRAY_OF : out FOOD_TYPE) do
 TRAY_OF := COOK;
end;
end loop;
end McD;

Task Body GONZO is

MY_TRAY : FOOD_TYPE;

procedure CONSUME (MY_TRAY : in FOOD_TYPE) is ...

begin

loop

McD.SERVE (MY_TRAY);

CONSUME (MY_TRAY);

end loop;

end GONZO;



```
Task Body Gonzo is
    NEW_TRAY : FOOD_TYPE;

    function COOK return FOOD_TYPE is
        ...
    end COOK;

begin
    loop
        NEW_TRAY := COOK;
        accept SERVE (TRAY_OF : out FOOD_TYPE ) do
            TRAY_OF := NEW_TRAY;
        end SERVE;
    end loop;
end GONZO;
```

```
loop  
    NEW_TRAY := COOK;  
    select  
        accept SERVE (TRAY_OF : out FOOD_TYPE) do...  
            TRAY_OF := NEW_TRAY;  
        end SERVE;  
    else  
        null;  
    end select;  
  
end loop;
```



```
loop  
    NEW_TRAY := COOK;  
    select  
        accept SERVE (TRAY_OF : out FOOD_TYPE) do...  
            TRAY_OF := NEW_TRAY;  
        end SERVE;  
        else  
            terminate;  
        end select;  
  
    end loop;
```



loop

NEW_TRAY := COOK;

select

accept SERVE (TRAY_OF : out FOOD_TYPE) do...

 TRAY_OF := NEW_TRAY;

end SERVE;

or

 delay 15 * MINUTES;

end select;

end loop;

```
loop  
    select  
        McD.SERVE(MY_ORDER); consume(MY_ORDER);  
    else  
        select  
            BK.SERVE(MY_ORDER); consume(MY_ORDER);  
        else  
            exit;  
        end select;  
    end select;  
  
end loop;
```



```
loop
    select
        McD.SERVE(MY_ORDER);  consume (MY_ORDER);
    or
        delay 10.0 * MINUTES;
    select
        BK.SERVE(MY_ORDER);  consume (MY_ORDER);
    or
        delay 5.0 * MINUTES;
        exit;
    end select;
end select;
end loop;
```

```
loop  
    select  
        MCD.SERVE (MY_ORDER);  
    or  
        BK.SERVE (MY_ORDER);  
    end select;  
  
    CONSUME;  
  
end loop;
```



loop

select

McD.SERVE (MY_ORDER);

or

BK.SERVE (MY_ORDER);

else

delay 10 * MINUTES;

exit;

end select;

consume;

end loop;



ADA TASKING

SCENARIO II

"No FREE LUNCH"

MCD TASK

SERVICE PROVIDED : Food

SERVICE REQUESTED: Money

GONZO TASK

SERVICE PROVIDED : Money

SERVICE REQUESTED: Food

```
Task McD is
    entry SERVE ( ORDER : out FOOD_TYPE;
                  COST : in MONEY_TYPE);
end McD;

TASK GONZO;
```

--OR

```
Task McD is
    entry SERVE ( ORDER : out FOOD_TYPE);
end McD;
```

```
Task GONZO is
    entry PAY ( COST      : in MONEY_TYPE;
                PAYMENT   : out MONEY_TYPE);
end GONZO;
```



```
Task Body McD is
  CASH_DRAWER : MONEY_TYPE;
  NEW_ORDER   : FOOD_TYPE;
  function COOK .....;
  function CALC_COST (ORDER : in FOOD_TYPE ) 
    return MONEY_TYPE is .....
begin
  loop
    NEW_ORDER := COOK;
    select
      accept SERVE(ORDER : out FOOD_TYPE) do
        ORDER := NEW_ORDER;
        COST := CALC_COST (NEW_ORDER);
        GONZO_PAY (COST, AMOUNT_PAID);
        CASH_DRAWER := CASH_DRAWER + AMOUNT_PAID;
      end SERVE;
      or
        delay 15.0 * MINUTES;
    end select;
  end loop;
end McD;
```



Task Body GONZO is

```
ACCOUNT_BALANCE : MONEY_TYPE;
MY_ORDER : FOOD_TYPE;
function GO_TO_WORK return MONEY_TYPE is.....
begin
  ACCOUNT_BALANCE := GO_TO_WORK + ACCOUNT_BALANCE;
loop
  McD.SERVE (MY_ORDER);
  accept PAY (COST : in MONEY_TYPE;
              PAYMENT : out MONEY_TYPE) do
    ACCOUNT_BALANCE := ACCOUNT_BALANCE -
                      COST ;
    PAYMENT := COST;
  end PAY;
end loop;
end GONZO;
```



ADA TASKING

SCENARIO II A

"NO WAIT FOR THE WAITERS"

MCD Task

SERVICE PROVIDED : FOOD
SERVICE REQUESTED: MONEY

GONZO Task

SERVICE PROVIDED : MONEY
SERVICE REQUESTED: FOOD

MANAGER Task

SERVICE PROVIDED : MAKE NEW WAITER
SERVICE REQUESTED: None

Task type McD is
entry SERVE....
end McD;

Task GONZO is
entry PAY....
end GONZO;

Task MANAGER;

Type CASHIER_POINTER is access McD;

Type REGISTER_TYPE is array (1..NO_REGS)
of CASHIER_POINTER;

THE_REGS : REGISTER_TYPE:=(others => new McD);

Task Body McD is

```
    ...
    ...
    ...
begin
loop
    NEW_ORDER := COOK;
    select
        accept SERVE.....
        ...
        end SERVE;
    or
        delay 2.0 * MINUTES;
        exit;
    end select;
end loop;
```

Task Body GONZO is

...

...

begin

...

...

--Now, GONZO has to search for the open
-- registers, and select the one with
-- the shortest line

...

...

THE_REGISTERS(MY_REGISTER).SERVE;

...

end GONZO;

```
Task Body MANAGER is
  ...
  ...
begin
  loop
    --The MANAGER will look at the queue lengths of
    -- the open registers, and, when necessary
    -- will open registers that are currently
    -- closed
    ...
    if ..... then
      THE_REGISTERS(CLOSED_REGISTER) := new McD;
      end if;
    end loop;
end MANAGER;
```

ADA TASKING

SCENARIO III

"A SUGAR CONE, PLEASE:

BR TASK
SERVICE PROVIDED : ICE CREAM
SERVICE REQUESTED: AN ORDER

SERVOMATIC TASK
SERVICE PROVIDED : A NUMBER

CUSTOMERS TASK
SERVICE PROVIDED : AN ORDER
SERVICE REQUESTED: ICE CREAM



```
task BR is
    entry SERVE (ICE_CREAM : out DESSERT_TYPE);
end BR;

task SERVOMATIC is
    entry TAKE (A_NUMBER : out SERVOMATIC_NUMBERS);
end SERVOMATIC;

task type CUSTOMER_TASK is
    entry REQUEST (ORDER : out ORDER_TYPE);
end CUSTOMER_TASK;

type CUSTOMER is access CUSTOMER_TASK;
CUSTOMERS : array (SERVOMATIC_NUMBERS) of CUSTOMER;
```



```
task body BR
    NEXT_CUSTOMER : SERUOMATIC_NUMBERS := SERUOMATIC_NUMBERS'last;
    CURRENT_ORDER : ORDER_TYPE;
    ICE_CREAM : DESSERT_TYPE;
    function MAKE (ORDER : in ORDER_TYPE) return DESSERT_TYPE is.....
begin
    loop
        begin
            NEXT_CUSTOMER := (NEXT_CUSTOMER + 1)
                mod SERUOMATIC_NUMBERS'last;
            CUSTOMERS(NEXT_CUSTOMER).REQUEST
                (CURRENT_ORDER);
            ICE_CREAM := MAKE(CURRENT_ORDER);
            accept SERVE(ICE_CREAM : out DESSERT_TYPE) do
                ICE_CREAM := BR.ICE_CREAM;
            end SERVE;
        exception
            when TASKING_ERROR => null;
                --customer not here
        end;
    end loop;
end;
```



```
task body SERUOMATIC is
    NEXT_NUMBER : SERUOMATIC_NUMBERS := SERUOMATIC_NUMBERS'first;
begin
    loop
        accept TAKE(A_NUMBER : out SERUOMATIC_NUMBERS) do
            A_NUMBER := NEXT_NUMBER;
        end TAKE;
        NEXT_NUMBER := (NEXT_NUMBER + 1) mod SERUOMATIC_NUMBERS'last;
    end loop;
end SERUOMATIC;
```

```
task body CUSTOMER_TASK is
    MY_ORDER : ORDER_TYPE := ... --some value;
    MY_DESSERT : DESSERT_TYPE;
begin
    accept REQUEST ( ORDER : out ORDER_TYPE) do
        ORDER := MY_ORDER;
    end REQUEST;
    BR.SERVE(MY_DESSERT);
    --eat the dessert, or do whatever
end;
```



ADA TASKING

SCENARIO IV

"LETS HIDE THE SPOOLER TASK"

PRINTER_PACKAGE

ACTION- "HIDES" THE PRINT SPOOLER
BY RENAMING TASK ENTRY

SPOOLER_TASK

SERVICE PROVIDED : VIRTUAL PRINT
SERVICE REQUESTED : PHYSICAL PRINT

PRINTER_Task

SERVICE PROVIDED : PHYSICAL PRINT
SERVICE REQUESTED: FILE NAME

```
Package PRINTER_PACKAGE is
  ...
  ...
  task SPOOLER is
    entry PRINT_FILE (NAME : in STRING;
                      PRIORITY : in NATURAL);
    entry PRINTER_READY;
  end SPOOLER;
  ...
  ...
  procedure PRINT ( NAME : in STRING;
                     PRIORITY : in NATURAL := 10)
    renames SPOOLER.PRINT_FILE;
  end PRINTER_PACKAGE;
```

Package Body PRINTER_PACKAGE is

```
  ...
  ...
  task PRINTER is
    entry PRINT_FILE (NAME : in STRING );
  end PRINTER;
  ...
  ...
end PRINTER_PACKAGE;
```

```
task body SPOOLER is
begin
  loop
    select
      accept PRINTER_READY do
        PRINTER.PRINT_FILE ( REMOVE (QUEUE) );
        -- Remove would determine the next job and
        -- send it to the actual printer
      end PRINTER_READY;
    else
      null;
    end select;
    select
      accept PRINT_FILE ( NAME : in STRING;
                           PRIORITY : NATURAL ) do
        INSERT ( NAME, PRIORITY );
        --put name on queue or queues according
        -- to priority
      end PRINT_FILE;
    else
      null;
    end select;
  end select;
  end loop;
end SPOOLER;
```



```
task body PRINTER is
begin
    loop
        SPOOLER.PRINTER_READY;
        accept PRINT FILE ( NAME : in STRING ) do
            if NAME'length /= 8 then .....
                --print the file
            else
                delay 10.0 * seconds;
            end if;

            end PRINT_FILE;
        end loop;
    end PRINTER;
```

```
with PRINTER_PACKAGE;

procedure MAIN is

    --
    --
    loop
        --process several files
        PRINTER_PACKAGE.PRINT (A_FILE, A_PRIORITY);
        --
    end loop;
end MAIN;
```

APPLICATIONS FOR TASKS

- CONCURRENT OPERATIONS
- ROUTING MESSAGES
- SHARED RESOURCE MANAGEMENT
- INTERRUPT HANDLING

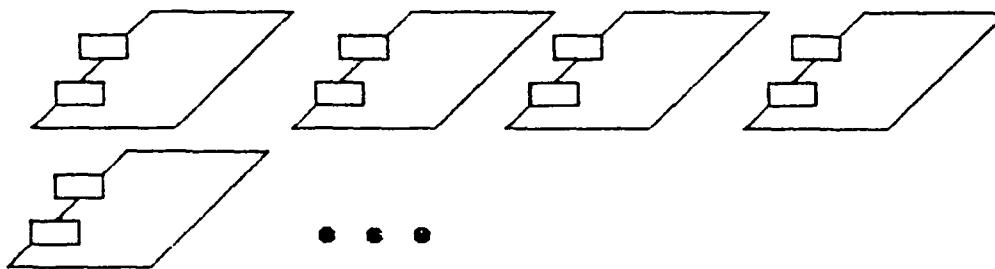
MATRIX MULTIPLICATION

$$\begin{bmatrix} 1 & 1 & 1 \\ 2 & 2 & 0 \end{bmatrix} * \begin{bmatrix} 2 \\ 1 \\ 1 \end{bmatrix} = \begin{bmatrix} 4 \\ 6 \end{bmatrix}$$

```
type ROW_OR_COL is array (integer range <>) of integer;
type PTR is access ROW_OR_COL;
```

```
task type PARTIAL is
    entry SEND (ROW, COL : ROW_OR_COL);
    entry RECEIVE (RESULT : out integer);
end PARTIAL;
```

MAIN



```
begin
    -- send row and col
    -- receive partial product
end
```

```
task body PARTIAL is
```

```
    PRODUCT : integer := 0;  
    ROW_PTR : PTR;  
    COL_PTR : PTR;
```

```
begin
```

```
    accept SEND (ROW,COL : ROW_OR_COL) do  
        ROW_PTR := new ROW_OR_COL'(ROW);  
        COL_PTR := new ROW_OR_COL'(COL);  
    end SEND;
```

```
    for I in ROW_PTR.all'range  
    loop  
        PRODUCT := PRODUCT +  
            ROW_PTR(I) * COL_PTR(I);  
    end loop;
```

```
    accept RECEIVE (RESULT : out integer) do  
        RESULT := PRODUCT;  
    end RECEIVE;
```

```
end PARTIAL;
```

```
procedure MAIN is

    COLS constant := 10;
    ROWS constant := 10;
    type MATRIX is array (1 .. ROWS) of
        ROW_OR_COL (1 .. COLS);

    MAT : MATRIX;
    VECTOR : ROW_OR_COL (1 .. COLS);
    FINAL : ROW_OR_COL (1 .. ROWS);

    ...

declare

    WORKER : array (1 .. ROWS) of PARTIAL; -- tasks

begin

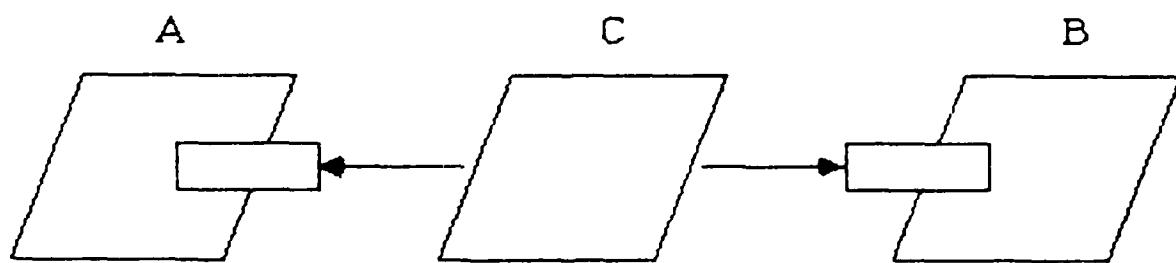
    for I in 1 .. ROWS
    loop
        WORKER(I).SEND(ROW => MAT(I),
                        COL => VECTOR);
    end loop;

    for I in 1 .. ROWS
    loop
        WORKER(I).RECEIVE (FINAL(I));
    end loop;

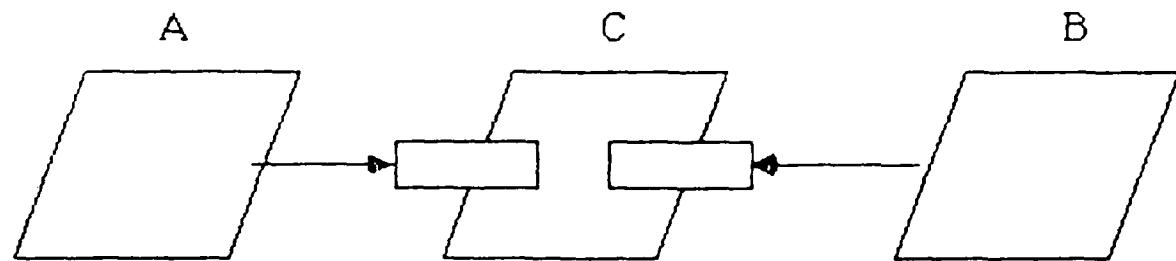
end;      -- block
```

ROUTING MESSAGES

- WRITE TASK SPECIFICATIONS TO SEND AN INTEGER FROM TASK A TO TASK B



- WRITE SPECIFICATIONS AND BODIES FOR THE FOLLOWING SYSTEM. TASK C WILL REPEATEDLY GET AN INTEGER FROM TASK A AND SEND IT ON TO TASK B



PRIORITY MESSAGES

type PRIORITY is (LOW, MEDIUM, HIGH),

task SWITCH is

```
entry SEND (PRIORITY)
      (M : in string);
```

end SWITCH;

task body SWITCH is

begin

loop

select

```
accept SEND(HIGH) do ... end SEND;
```

or

```
when SEND(HIGH)'count = 0 =>
accept SEND(MEDIUM) do ... end SEND;
```

or

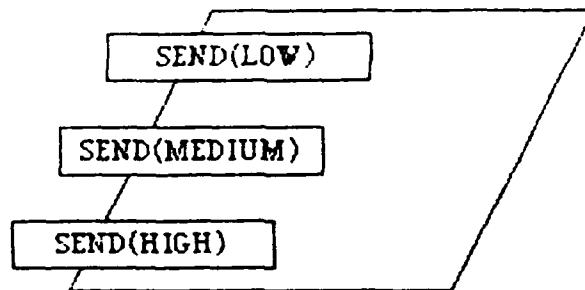
```
when SEND(HIGH)'count = 0 and
SEND(MEDIUM)'count = 0 =>
accept SEND (LOW) ... end SEND;
```

end select;

end loop;

end SWITCH;

SWITCH



A Synchronizing Buffer

```
task SYNCHRONIZER is
    entry PUT (ITEM : in SOME_TYPE);
    entry GET (ITEM : out SOME_TYPE);
end SYNCHRONIZER;

task body SYNCHRONIZER is
    SPOT : SOME_TYPE;

begin
    loop
        accept PUT (ITEM : in SOME_TYPE) do
            SPOT := ITEM;
        end PUT;

        accept GET (ITEM : out SOME_TYPE) do
            ITEM := SPOT;
        end GET;

    end loop;
end SYNCHRONIZER;
```

CONTROLLING RESOURCES

- SEVERAL CONCERNS ARE PRESENT WHEN DEALING WITH PARALLELISM THAT ARE NOT PRESENT WHEN DEALING IN A PURELY SEQUENTIAL MODE
- IT IS IMPORTANT TO BE ABLE TO ASSURE THAT A VALUE IS NOT BEING CHANGED BY ONE USER AT THE PRECISE MOMENT THAT IT IS BEING REFERENCED BY ANOTHER USER
- Ada PROVIDES A PRAGMA 'SHARED' WHICH CAN HELP

```
INDEX : integer;  
pragma SHARED(INDEX);
```

- ENFORCES MUTUALLY EXCLUSIVE ACCESS
- AVAILABLE FOR SCALAR AND ACCESS TYPES ONLY

SEMAPHORES

```
task SEMAPHORE is
    entry SEIZE;
    entry RELEASE;
end SEMAPHORE;
```

```
task body SEMAPHORE is
    IN_USE : boolean := false;
begin
    loop
        select
```

```
when not IN_USE =>
    accept SEIZE do
        IN_USE := true;
    end SEIZE;
```

or

```
when IN_USE =>
    accept RELEASE do
        IN_USE := false;
    end RELEASE;
```

```
end select;
end loop;
end SEMAPHORE;
```

ENCAPSULATING A DATA ITEM

```
task PROTECTED is
    entry SET (OBJ : in integer);
    entry GET (OBJ : out integer);
end PROTECTED;
```

```
task body PROTECTED is
```

```
    LOCAL : integer; 
begin
    loop
        select
```

```
            accept SET (OBJ : in integer) do
                LOCAL := OBJ;
            end SET;
```

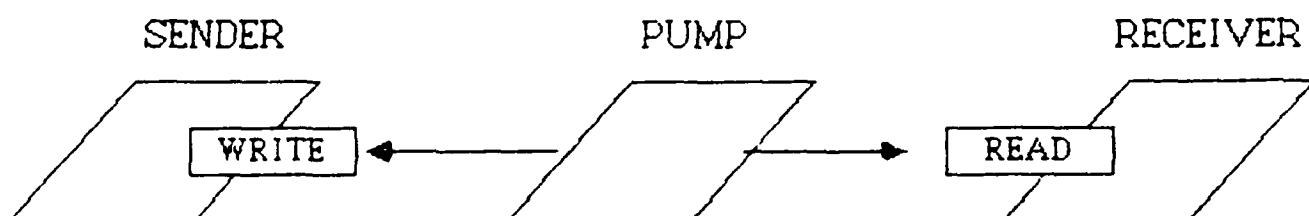
or

```
            accept GET (OBJ : out integer) do
                OBJ := LOCAL;
            end GET;
```

```
        end select;
    end loop;
end PROTECTED;
```

PUMPING TASK

```
task PUMP;  
  
task SENDER is  
    entry WRITE (ITEM : out SOME_TYPE);  
end SENDER;  
  
task RECEIVER is  
    entry READ (ITEM : in SOME_TYPE);  
end RECEIVER;  
  
task body PUMP is  
    THE_ITEM : SOME_TYPE;  
begin  
    loop  
        SENDER.WRITE(THE_ITEM);  
        RECEIVER.READ(THE_ITEM);  
    end loop;  
end PUMP;  
  
task body SENDER is separate;  
task body RECEIVER is separate;
```



HARDWARE INTERRUPTS

- FOR ARCHITECTURES THAT 'JUMP' TO A CERTAIN HARDWARE ADDRESS UPON RECEIPT OF AN INTERRUPT
- A TASK ENTRY IS ASSOCIATED WITH THE ADDRESS
- PRIORITY IS HIGHER THAN ANY USER-DEFINED

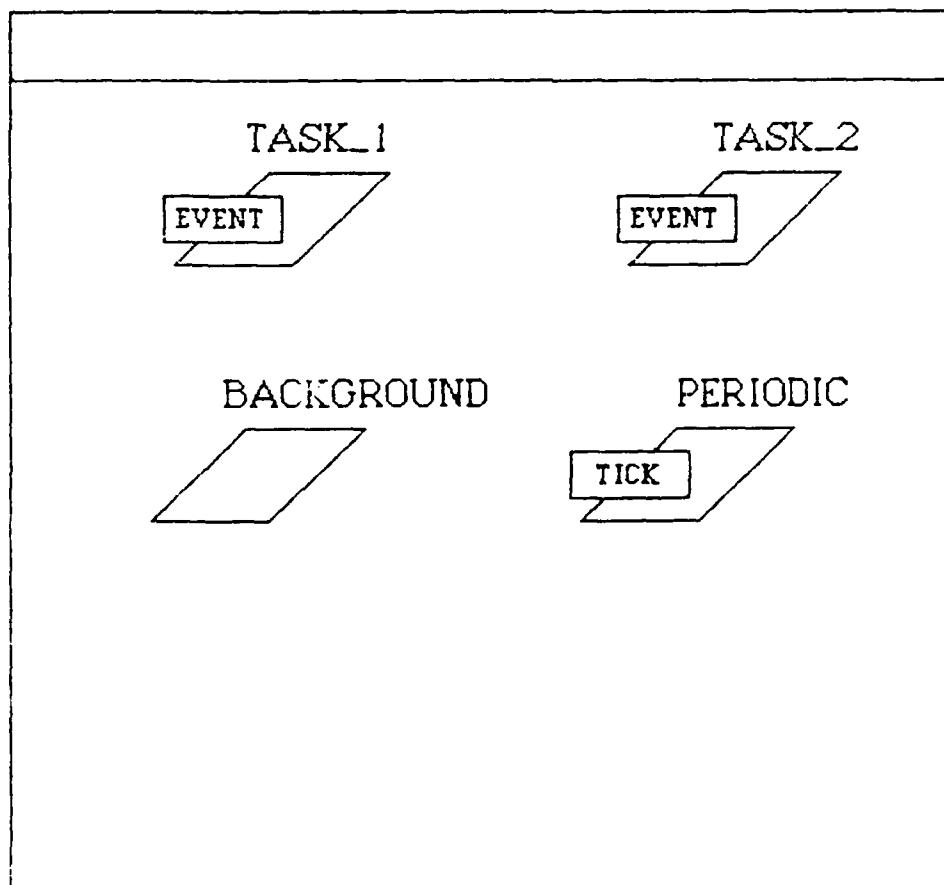
```
task INTERRUPT_HANDLER is
    entry DONE;
        for DONE use at 16#40#;
end INTERRUPT_HANDLER;
```

```
task body INTERRUPT_HANDLER is
begin
    accept DONE do
        ...
    end DONE;
end INTERRUPT_HANDLER;
```

EVENT DRIVEN SYSTEMS W/BACKGROUND

- A cyclic executive might deal with several levels of processing
 - Event driven processing (high priority, perhaps interrupt handling)
 - Periodic (cyclic) processing
 - Background processing (low priority)

EXECUTIVE



```
procedure EXECUTIVE is

    task TASK_1 is
        pragma PRIORITY (10);
        entry EVENT;
    end TASK_1;

    task TASK_2 is
        entry EVENT;
        for EVENT use at 16#110#;
    end TASK_2;

    task BACKGROUND is
        pragma PRIORITY (0);
    end BACKGROUND;

    task PERIODIC is
        pragma PRIORITY (5);
        entry TICK;           -- one tick per cycle
    end PERIODIC;

    task body PERIODIC is
        ...
        begin
            loop
                accept TICK;
                ... -- process a frame
            end loop;
        end PERIODIC;

        -- bodies (or stubs) of other tasks go here

    end EXECUTIVE;
```

Second Annual ASEET Symposium

Tutorial on Ada[®] Exceptions

by
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Air Force Institute of Technology (AFIT)
and
Arizona State University (ASU)

9 June 1987

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Outline

=> Overview

- Naming an exception
- Creating an exception handler
- Raising an exception
- Handling exceptions
- Turning off exception checking
- Tasking exceptions
- More examples

Overview

- What is an exception
- Ada exceptions
- Comparison
 - the American way
 - using exceptions

What Is an Exception

- A run time error
- An unusual or unexpected condition
- A condition requiring special attention
- Other than normal processing

Ada Exceptions

- An exception has a name
 - may be predefined
 - may be declared
- The exception is raised
 - may be raised implicitly by run time system
 - may be raised explicitly by **raise** statement
- The exception is handled
 - exception handler may be placed in any **frame**
 - exception propagates until handler is found
 - if no handler anywhere, process aborts

The American Way

```
package Stack_Package is

    type Stack_Type is limited private;

    procedure Push (Stack : in out Stack_Type;
                    Element : in Element_Type;
                    Overflow_Flag : out boolean);
    ...

end Stack_Package;

with Text_IO;
with Stack_Package; use Stack_Package;
procedure Flag_Waving is
    ...
    Stack : Stack_Type;
    Element : Element_Type;
    Flag : boolean;
begin
    ...
    Push (Stack, Element, Flag);
    if Flag then
        Text_IO.Put ("Stack overflow");
        ...
    end if;
    ...
end Flag_Waving;
```

Using Exceptions

```
package Stack_Package is

    type Stack_Type is limited private;
    Stack_Overflow,
    Stack_Underflow : exception;

    procedure Push (Stack : in out Stack_Type;
                    Element : in Element_Type);
                    -- may raise Stack_Overflow
    ...
end Stack_Package;
```

```
with Text_IO;
with Stack_Package; use Stack_Package;
procedure More_Natural is
    ...
    Stack : Stack_Type;
    Element : Element_Type;
begin
    ...
    Push (Stack, Element);
    ...
exception
    when Stack_Overflow =>
        Text_IO.Put ("Stack overflow");
    ...
end More_Natural;
```

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Naming an Exception

- Predefined exceptions
- Declaring exceptions
- I/O exceptions

Predefined Exceptions

- In package STANDARD (also see chap 11 of LRM)
- CONSTRAINT_ERROR
 - violation of range, index, or discriminant constraint...
- NUMERIC_ERROR
 - execution of a predefined numeric operation cannot deliver a correct result
- PROGRAM_ERROR
 - attempt to access a program unit which has not yet been elaborated...
- STORAGE_ERROR
 - storage allocation is exceeded...
- TASKING_ERROR
 - exception arising during intertask communication

Declaring Exceptions

```
exception_declaration ::= identifier_list : exception;
```

- Exception may be declared anywhere an object declaration is appropriate
- However, exception is not an object
 - may not be used as subprogram parameter, record or array component
 - has same scope as an object, but its effect may extend beyond its scope

Example:

```
procedure Calculation is  
    Singular : exception;  
    Overflow, Underflow : exception;  
  
begin  
    ...  
end Calculation;
```

I/O Exceptions

- Exceptions relating to file processing
- In predefined library unit IO_EXCEPTIONS
(also see chap 14 of LRM)
- TEXT_IO, DIRECT_IO, and SEQUENTIAL_IO with it

package IO_EXCEPTIONS is

| | |
|---------------------------|---|
| NAME_ERROR : exception; | |
| USE_ERROR : exception; | --attempt to use --invalid operation |
| STATUS_ERROR : exception; | |
| MODE_ERROR : exception; | |
| DEVICE_ERROR : exception; | |
| END_ERROR : exception; | --attempt to read --beyond end of file |
| DATA_ERROR : exception; | --attempt to input --wrong type |
| LAYOUT_ERROR : exception; | --for text processing |

end IO_EXCEPTIONS;

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Creating an Exception Handler

- Defining an exception handler
- Restrictions
- Handler example

Defining an Exception Handler

- Exception condition is "caught" and "handled" by an exception handler
- Exception handler may appear at the end of any frame (block, subprogram, package or task body)

```
begin
  ...
exception
  -- exception handler(s)
end;
```

- Form similar to case statement

```
exception_handler ::=  
  when exception_choice { | exception_choice} =>  
    sequence_of_statements
```

```
exception_choice ::= exception_name | others
```

Restrictions

- Exception handlers must be at the end of a frame
- Nothing but exception handlers may lie between **exception** and **end** of frame
- A handler may name any visible exception declared or predefined
- A handler includes a sequence of statements
 - response to exception condition
- A handler for **others** may be used
 - must be the last handler in the frame
 - handles all exceptions not listed in previous handlers of the frame
(including those not in scope of visibility)
 - can be the only handler in the frame

Handler Example

```
procedure Whatever is
    Problem_Condition : exception;
begin
    ...
exception
    when Problem_Condition =>
        Fix_It;
    when CONSTRAINT_ERROR =>
        Report_It;
    when others =>
        Punt;
end Whatever;
```

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Raising an Exception

- How exceptions are raised
- Effects of raising an exception
- Raising example

How Exceptions are Raised

- Implicitly by run time system
 - predefined exceptions
- Explicitly by **raise** statement

raise_statement ::= **raise** [exception_name];

- the name of the exception must be visible at the point of the raise statement
- a raise statement without an exception name is allowed only within an exception handler

Effects of Raising an Exception

- Control transfers to exception handler at end of frame
(if one exists)
- Exception is lowered
- Sequence of statements in exception hander is executed
- Control passes to end of frame
- If frame does not contain an appropriate exception handler,
the exception is propagated

Raising Example

```
procedure Whatever is

    Problem_Condition : exception;
    Real_Bad_Condition : exception;

begin
    ...
    if Problem_Arises then
        raise Problem_Condition;
    end if;
    ...
    if Serious_Problem then
        raise Real_Bad_Condition;
    end if;
    ...
exception
    when Problem_Condition =>
        Fix_It;

    when CONSTRAINT_ERROR =>
        Report_It;

    when others =>
        Punt;
end Whatever;
```

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Handling Exceptions

- How exception handling can be useful
- Which exception handler is used
- Sequence of statements in exception handler
- Propagation
- Propagation example

How Exception Handling Can Be Useful

- Normal processing could continue if
 - cause of exception condition can be "repaired"
 - alternative approach can be used
 - operation can be retried
- Degraded processing could be better than termination
 - for example, safety-critical systems
- If termination is necessary, "clean-up" can be done first

Which Exception Handler Is Used

- If exception is raised during normal execution, system looks for an exception handler at the end of the frame in which the exception occurred
- If exception is raised during elaboration of the declarative part of a frame
 - elaboration is abandoned and control goes to the end of the frame with the exception still raised
 - exception part of the frame is not searched for an appropriate handler
 - effectively, the calling unit will be searched for an appropriate handler
 - if elaboration of library unit, program execution is abandoned
 - all library units are elaborated with the main program
- If exception is raised in exception handler
 - handler may contain block(s) with handler(s)
 - if not handled locally within handler, control goes to end of frame with exception raised

Sequence of Statements in Exception Handler

- Handler completes the execution of the frame
 - handler for a **function** should usually contain a **return** statement
- Statements can be of arbitrary complexity
 - can use most any language construct that makes sense in that context
 - cannot use **goto** statement to transfer into a handler
 - if handler is in a block inside a loop, could use **exit** statement
- Handler at end of package body applies only to package initialization

Propagation

- Occurs if no handler exists in frame where exception is raised
- Also occurs if **raise** statement is used in handler
- Exception is propagated dynamically
 - propagates from subprogram to unit calling it
(not necessarily unit containing its declaration)
 - this can result in propagation outside its scope
- Propagation continues until
 - an appropriate handler is found
 - exception propagates to main program (still with no handler) and program execution is abandoned

Propagation Example

```
procedure Do_Nothing is
-----
    procedure Has_It is
        Some_Problem : exception;
    begin
        ...
        raise Some_Problem;
        ...
    exception
        when Some_Problem =>
            Clean_Up;
            raise;
    end Has_It;
-----
    procedure Calls_It is
    begin
        ...
        Has_It;
        ...
    end Calls_It;
-----
begin -- Do_Nothing
    ...
    Calls_It;
    ...
exception
    when others => Fix_Everything;
end Do_Nothing;
```

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Turning Off Exception Checking

- Overhead vs efficiency
- Pragma SUPPRESS
- Check identifiers

Overhead vs Efficiency

- Exception checking imposes run time overhead
 - interactive applications will never notice
 - real-time applications have legitimate concerns but must not sacrifice system safety
- When efficiency counts
 - first and foremost, make program work
 - be sure possible problems are covered by exception handlers
 - check if efficient enough - stop if it is
 - if not, study execution profile
 - eliminate bottlenecks
 - improve algorithm
 - avoid "cute" tricks
 - check if efficient enough - stop if it is
 - if not, trade-offs may be necessary
 - some exception checks may be expendable since debugging is done
 - however, every suppressed check poses new possibilities for problems
 - must re-examine possible problems
 - must re-examine exception handlers
 - always keep in mind
 - problems will happen
 - critical applications must be able to deal with these problems

Moral

Improving the algorithm is far better - and easier in
the long run - than suppressing checks

Pragma SUPPRESS

- Only allowed immediately within a declarative part or immediately within a package specification

pragma SUPPRESS (identifier [, [ON =>] name]);

- identifier is that of the check to be omitted
(next slide lists identifiers)
- name is that of an object, type, or unit for which the check is to be suppressed
 - if no name is given, it applies to the remaining declarative region
- An implementation is free to ignore the suppress directive for any check which may be impossible or too costly to suppress

Example:

pragma SUPPRESS (INDEX_CHECK, ON => Index);

Check Identifiers

- These identifiers are explained in more detail in chap 11 of the LRM
- Check identifiers for suppression of CONSTRAINT_ERROR checks

ACCESS_CHECK
DISCRIMINANT_CHECK
INDEX_CHECK
LENGTH_CHECK
RANGE_CHECK

- Check identifiers for suppression of NUMERIC_ERROR checks

DIVISION_CHECK
OVERFLOW_CHECK

- Check identifier for suppression of PROGRAM_ERROR checks

ELABORATION_CHECK

- Check identifier for suppression of STORAGE_ERROR check

STORAGE_CHECK

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Tasking Exceptions

- Exception handling is trickier for tasks
- Exceptions during task rendezvous
- Tasking example

Exception Handling Is Trickier for Tasks

- Rules are not really different, just more involved
 - local exceptions handled the same within frames

If exception is raised

- during elaboration of task declarations
 - the exception TASKING_ERROR will be raised at the point of task activation
 - the task will be marked completed
- during execution of task body (and not resolved there)
 - task is completed
 - exception is not propagated
- during task rendezvous
 - this is the really tricky part

Exceptions During Task Rendezvous

- If the **called** task terminates abnormally
 - exception TASKING_ERROR is raised in **calling** task at the point of the entry call
- If the **calling** task terminates abnormally
 - no exception propagates to the **called** task
- If an exception is raised in **called** task within an **accept** (and not handled there locally)
 - the same exception is raised in the **calling** task at the point of the entry call
 - (even if exception is later handled outside of the accept in the called task)
- If an entry call is made for entry of a task that becomes completed before accepting the entry
 - exception TASKING_ERROR is raised in **calling** task at the point of the entry call

Tasking Example

```
procedure Critical_Code is

    Failure : exception;
    -----
    task Monitor is
        entry Do_Something;
    end Monitor;
    task body Monitor is
        ...
        begin
            accept Do_Something do
                ...
                raise Failure;
                ...
            end Do_Something;
            ...
        exception -- exception handled here
            when Failure =>
                Termination_Message;
        end Monitor;
        -----
        begin -- Critical_Code
            ...
            Monitor.Do_Something;
            ...
        exception -- same exception will be handled here
            when Failure =>
                Critical_Problem_Message;
        end Critical_Code;
```

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=> **More examples**

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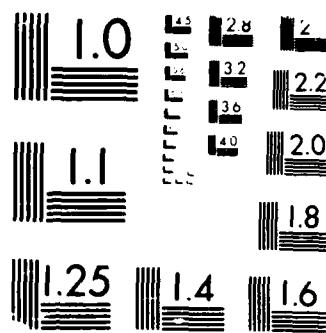
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Interactive Data Input

```
with Text_io; use Text_io;
procedure Get_Input (Number : out integer) is

    type Input_Type is integer range 0..100;
    package Int_io is new Integer_io (Input_Type);
    In_Number : Input_Type;

begin -- Get_Input

    loop      -- to try again after incorrect input

        begin -- inner block to hold exception handler

            put ("Enter a number 0 to 100");
            Int_io.get (In_Number);
            Number := In_Number;
            exit; -- to exit loop after correct input

        exception
            when DATA_ERROR | CONSTRAINT_ERROR =>
                put ("Try again, fat fingers!");
                Skip_Line; -- must clear buffer

        end; -- inner block

    end loop;

end Get_Input;
```

Propagating Exception Out of Scope and Back In

```
declare
    package Container is
        procedure Has_Handler;
        procedure Raises_Exception;
    end Container;
    -----
    procedure Not_in_Package is
    begin
        Container.Raises_Exception;
    exception
        when others => raise;
    end Not_in_Package;
    -----
    package body Container is
        Crazy : exception;
        procedure Has_Handler is
        begin
            Not_in_Package;
        exception
            when Crazy => Tell_Everyone;
        end Has_Handler;
        procedure Raises_Exception is
        begin
            raise Crazy;
        end Raises_Exception;
    end Container;
begin
    Container.Has_Handler;
end;
```

Keeping a Task Alive

```
task Monitor is
    entry Do_Something;
end Monitor;

task body Monitor is
begin
    loop      -- for never-ending repetition
        ...
        select
            accept Do_Something do
                begin -- block for exception handler
                    ...
                    raise Failure;
                    ...
                    exception
                        when Failure => Recover;
                    end; -- block
                end Do_Something; -- exception must be
                                    -- lowered before exiting
                ...
            end select;
            ...
        end loop;

    exception
        when others =>
            Termination_Message;
    end Monitor;
```

END

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